

**A computationally
universal phase
of
quantum
matter**

Robert Raussendorf, UBC

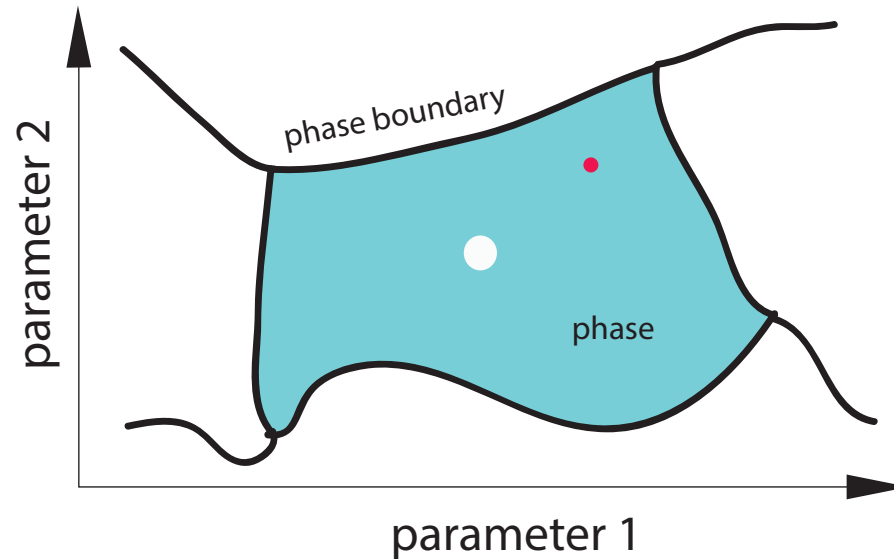
joint work with D.-S. Wang, D.T. Stephen, C. Okay, and H.P. Nautrup

The liquid phase of water



A quantum phase of spins in 2D

... which supports universal quantum computation

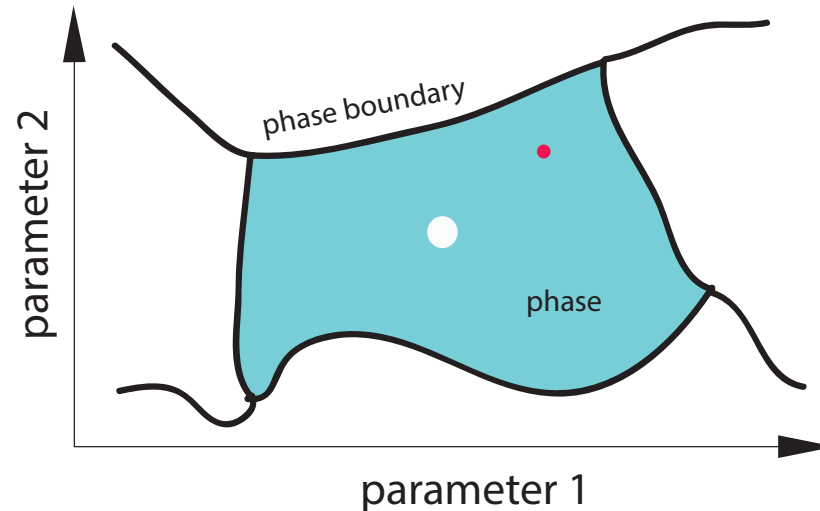


We consider:

- Phases of unique ground states of spin Hamiltonians, at $T = 0$,
- In the presence of symmetry,
- In spatial dimension 2 (a lattice of spin 1/2 particles)

A quantum phase of spins in 2D

... which supports universal quantum computation



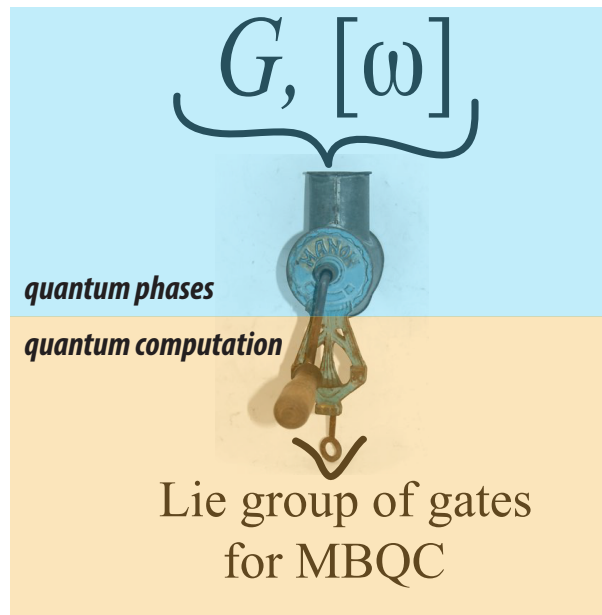
We show: for measurement-based quantum computation,

- There exists a quantum phase of matter which is universal for quantum computation
- The computational power is uniform across the phase.

Outline

1. “Computational phases of quantum matter”:
 - Our motivation
 - Background: SPT & MBQC
 - A short history of the question
2. A computationally universal phase of matter in 2D

Motivation #1: MBQC and symmetry



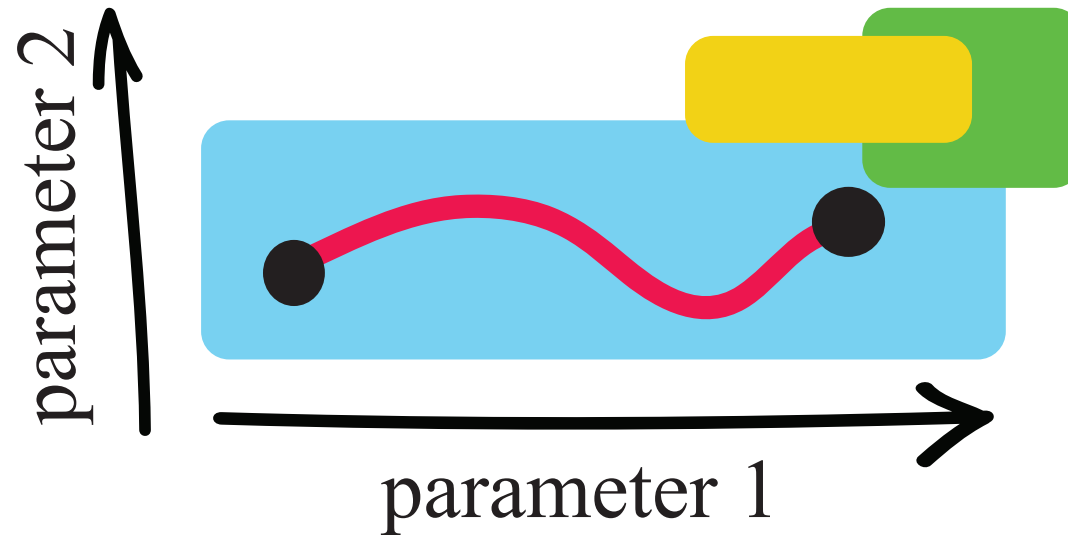
Can MBQC schemes be classified by symmetry, in a similar way as, say, elementary particles can?

If so, does this have a bearing on quantum algorithms?

Part I:

Background on SPT & MBQC

Symmetry-protected topological order

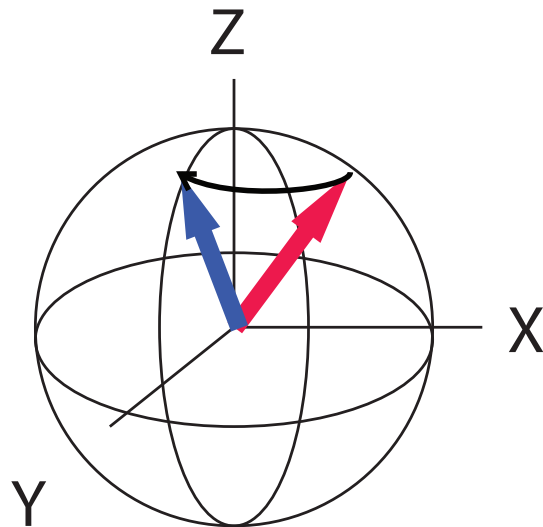


Two points in parameter space lie in the same SPT phase iff they can be connected by a path of Hamiltonians such that

1. At every point on the path, the corresponding Hamiltonian is invariant under G .
2. Along the path the energy gap never closes.

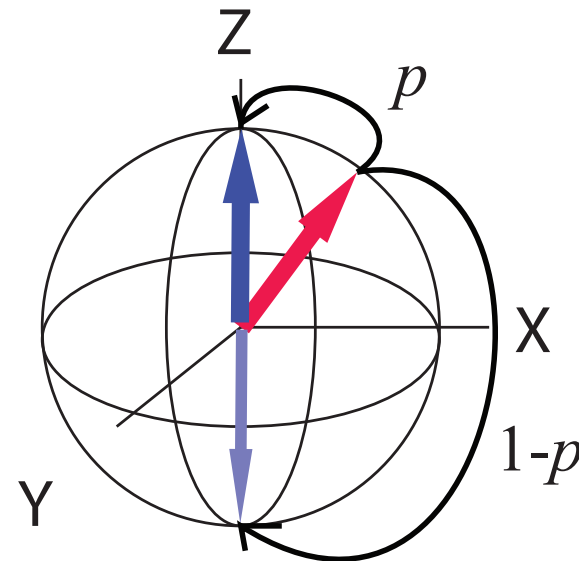
Measurement-based quantum computation

Unitary transformation



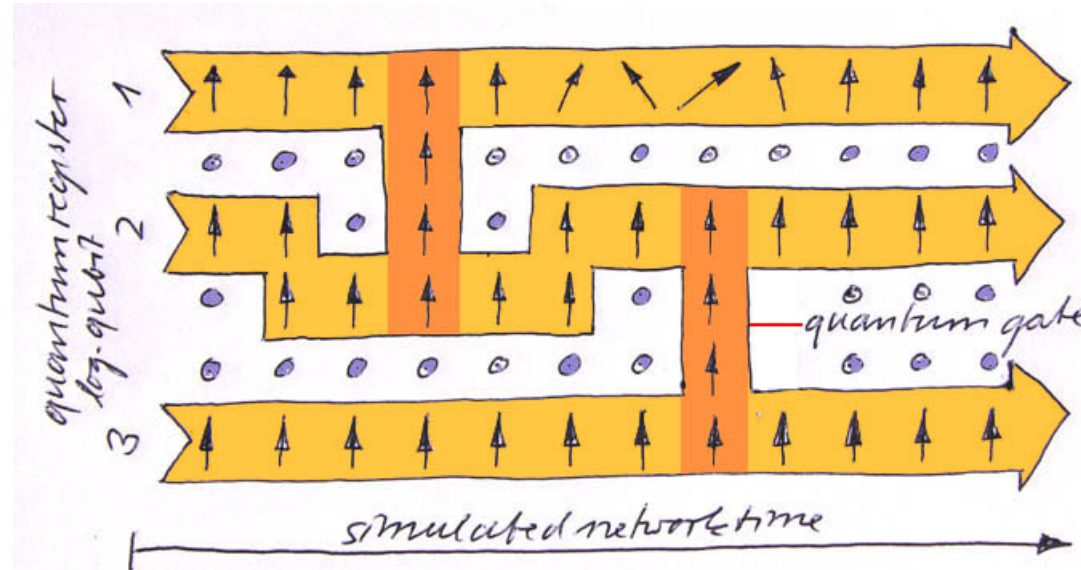
deterministic,
reversible

Projective measurement



probabilistic,
irreversible

Measurement-based quantum computation



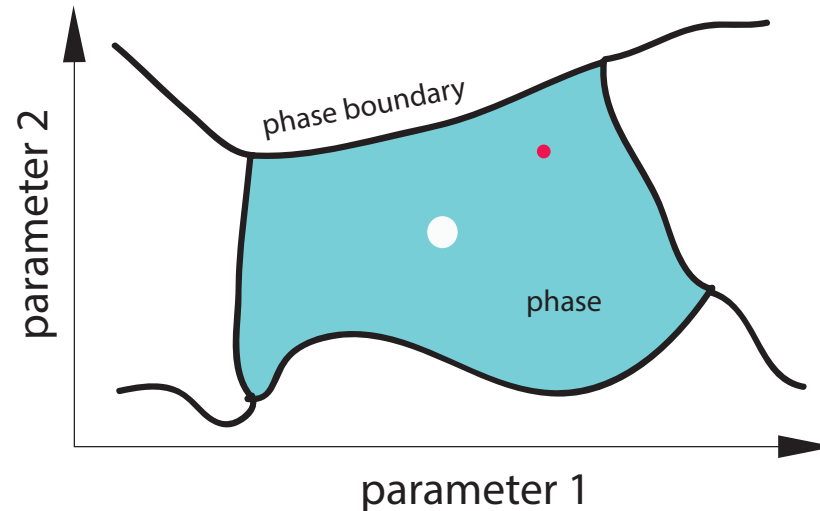
measurement of Z (\odot), X (\uparrow), $\cos \alpha X + \sin \alpha Y$ (\nearrow)

- Information written onto the resource state, processed and read out by one-qubit measurements only.
- Universal computational resources exist: cluster state, AKLT state.

R. Raussendorf, H.-J. Briegel, Physical Review Letters 86, 5188 (2001).

A quantum phase of spins in 2D

... which supports universal quantum computation



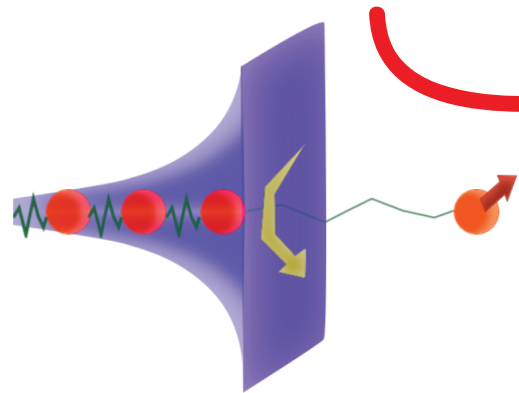
We show: for measurement-based quantum computation,

- There exists a quantum phase of matter which is universal for quantum computation
- The computational power is uniform across the phase.

*A short history of
“computational phases of quantum matter”*

1. Symmetry protects computation

we observe low-maintenance features of the ground-code MQC in that this computation is doable without an exact (classical) description of the resource ground state as well as without an initialization to a pure state. It

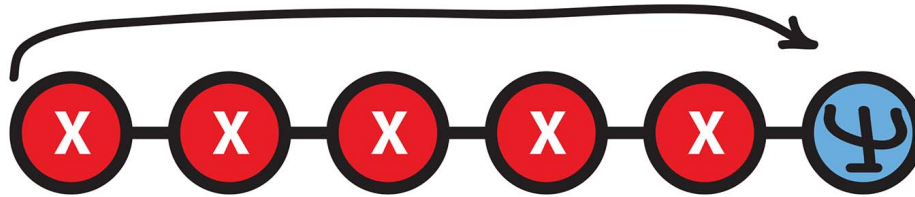


It turns out these features are deeply intertwined with the physics of the 1D Haldane phase (cf. Fig. 1), that is well characterized as the symmetry-protected topological order in a modern perspective [6, 7]. We believe our approach must bring the study of MQC, conventionally based on the analysis of the model entangled states (e.g., [1, 8, 9]), much closer to the condensed matter physics, which is aimed to describe characteristic physics based on the Hamiltonian.



A. Miyake, Phys. Rev. Lett. 105, 040501 (2010).

2. Symmetry-protected wire in MBQC

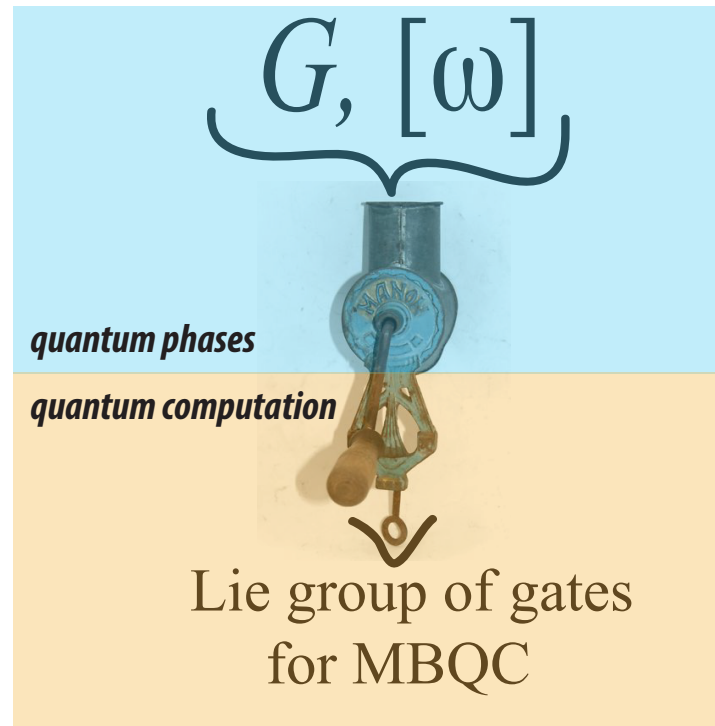


- Computational wire persists throughout symmetry-protected phases in 1D.
- Imports group cohomology from the classification of SPT phases.

D.V. Else, I. Schwartz, S.D. Bartlett and A.C. Doherty, PRL 108 (2012).

F. Pollmann *et al.*, PRB B 81, 064439 (2010); N. Schuch, D. Perez-Garcia, and I. Cirac, PRB 84, 165139 (2011); X. Chen, Z.-C. Gu, and X.-G. Wen, PRB 83, 035107 (2011).

3. The SPT \Rightarrow MBQC meat grinder



Hints at the classification of MBQC schemes by symmetry.

J. Miller and A. Miyake, Phys. Rev. Lett. 114, 120506 (2015) [first 1D comp. phase].

A. Prakash and T.-C. Wei, Phys. Rev. A (2016).

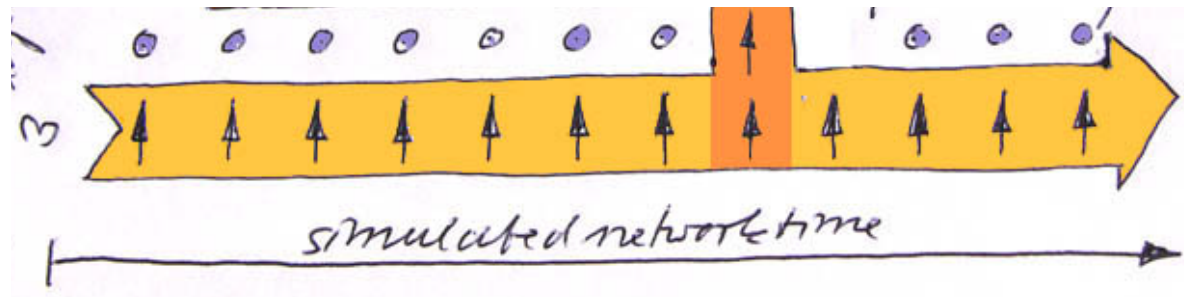
RR, A.Prakash, D.-S. Wang, T.-C.Wei, D.T. Stephen, Phys. Rev. A (2017).

Inspection

The above waypoints are about 1D systems.

1D is not sufficient for universal MBQC

here is why:



- MBQC in spatial dimension D maps to the circuit model in dimension $D - 1$

⇒ Require $D \geq 2$ for universality.

*Are there
computationally universal
quantum phases
in two dimensions?*

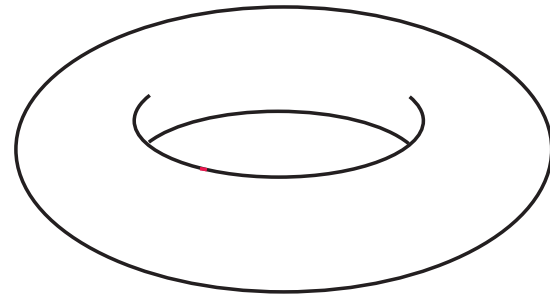
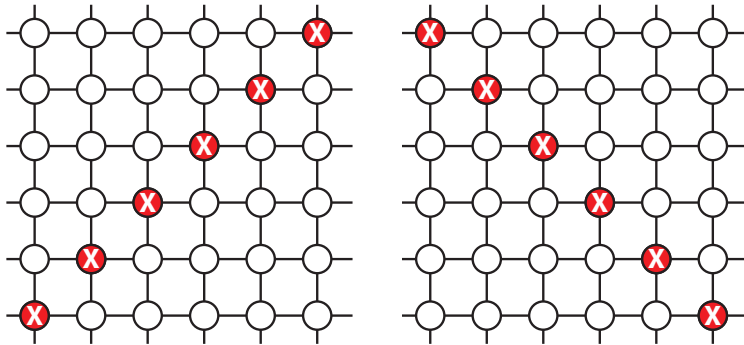
This talk describes one.

Part II:

A computationally universal SPT phase in 2D

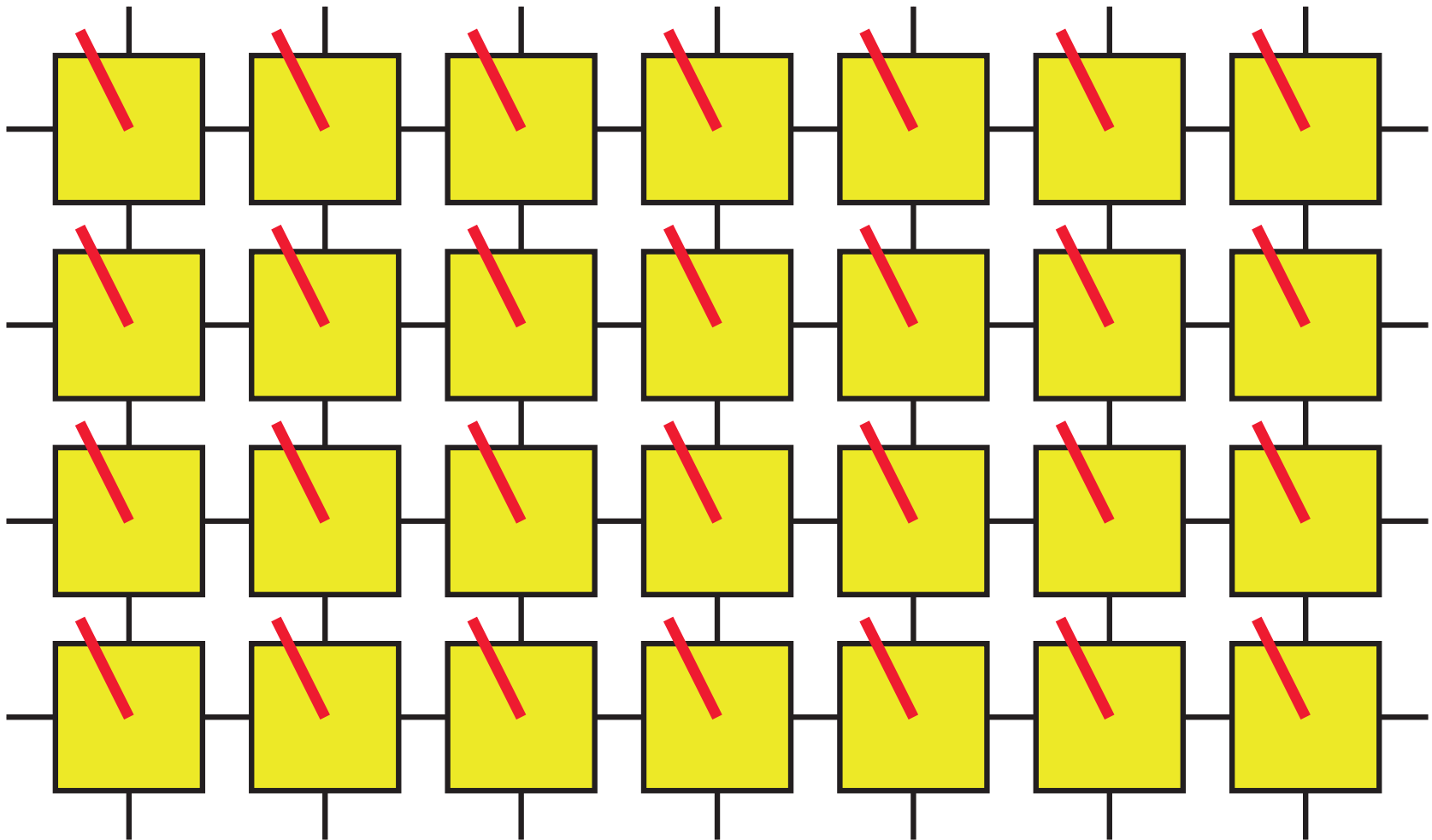
Description of the 2D phase & result

- The symmetries of the phase are



- The 2D cluster state is inside the phase

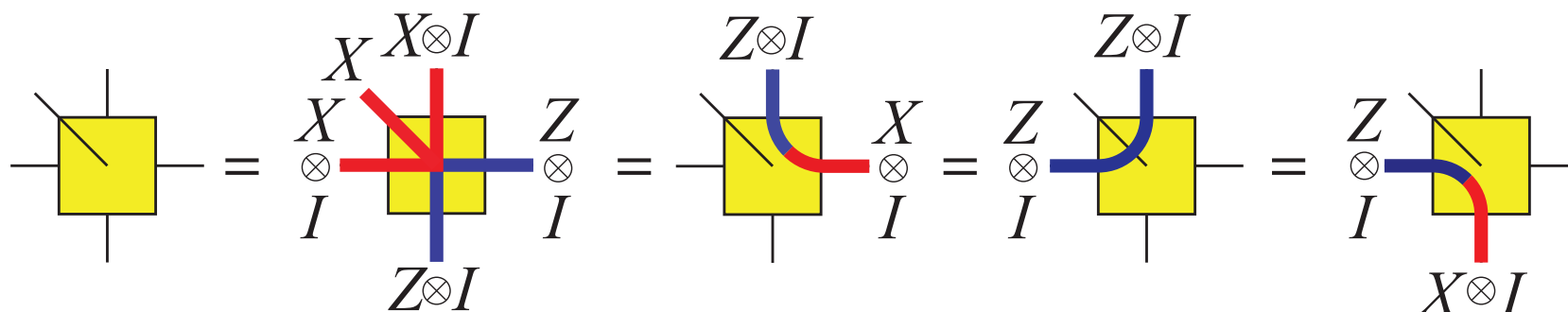
Result. For a spin-1/2 lattice on a torus with circumferences n and Nn , with n even, all ground states in the 2D cluster phase, except a possible set of measure zero, are universal resources for measurement-based quantum computation on $n/2$ logical qubits.



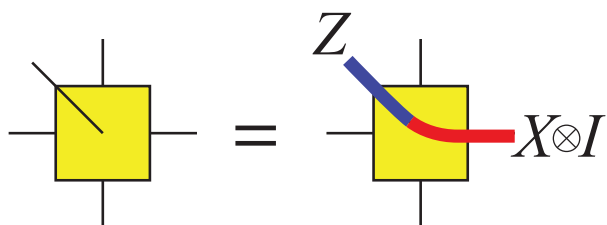
Consider MBQC resource states as tensor networks

Cluster-like states

... have PEPS tensors with the following symmetries



The cluster states have the additional symmetry



(We do not require the latter symmetry for cluster-like states)

Splitting the problem into halves

Part A:

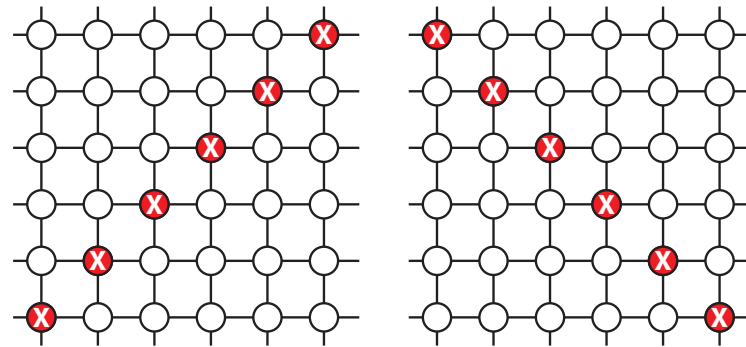
Lemma 1. All states in the 2D cluster phase are cluster-like.

Part B:

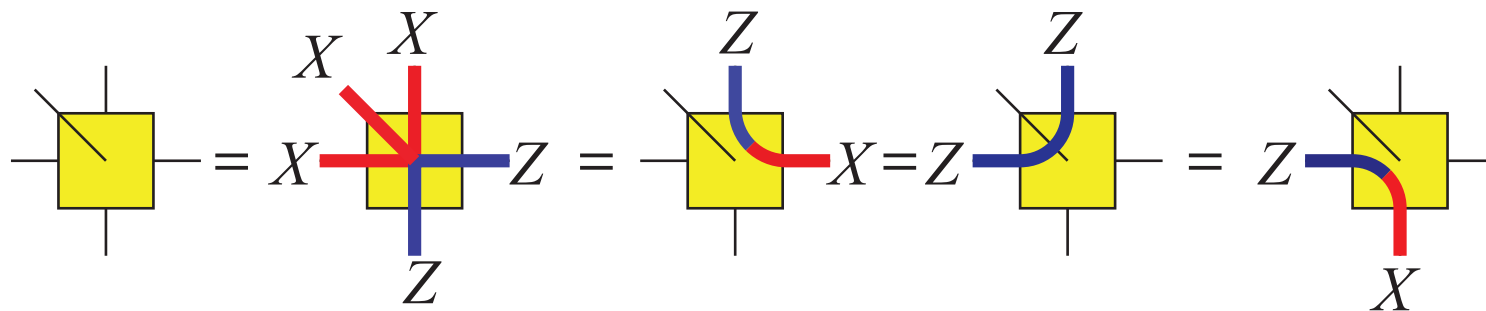
Lemma 2. All cluster-like states, except a set of measure zero, are universal for MBQC.

Part A: PEPS tensor symmetries

The physical symmetries

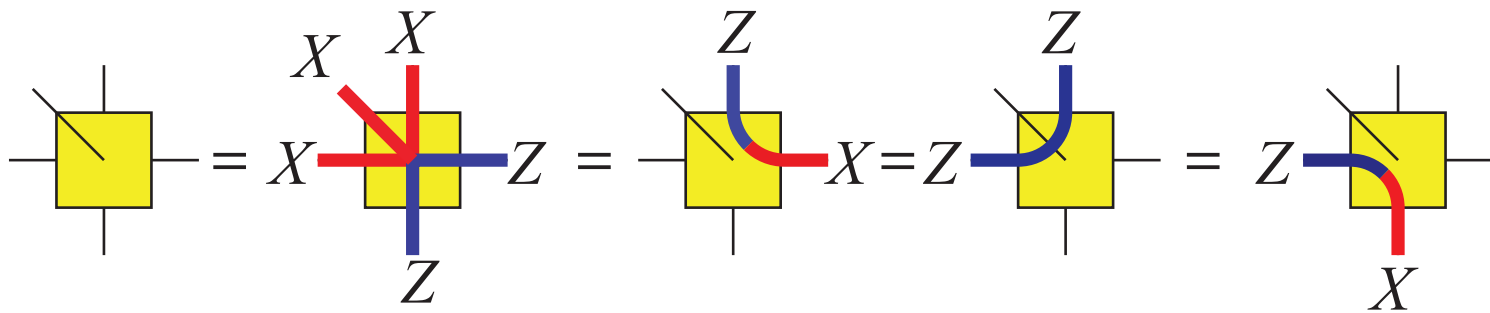


in the 2D cluster phase imply the local PEPS tensor symmetries,



Part B: Symmetry Lego

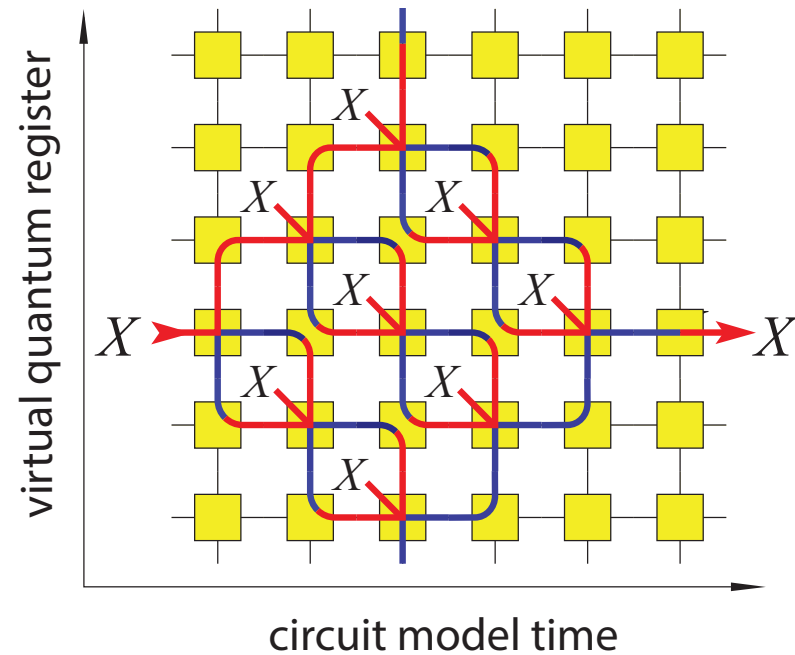
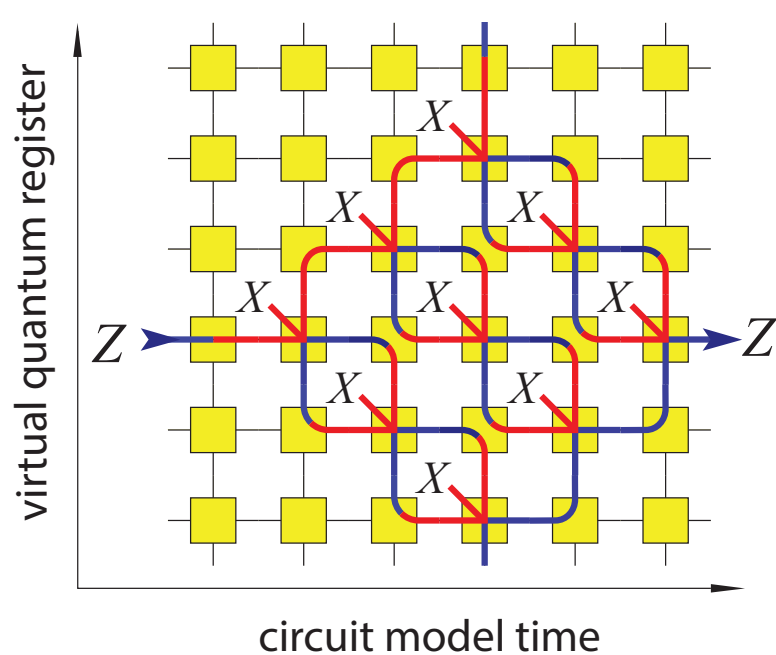
Now weave the PEPS tensor symmetries



into larger patterns.

B: Cluster-like \Rightarrow universal

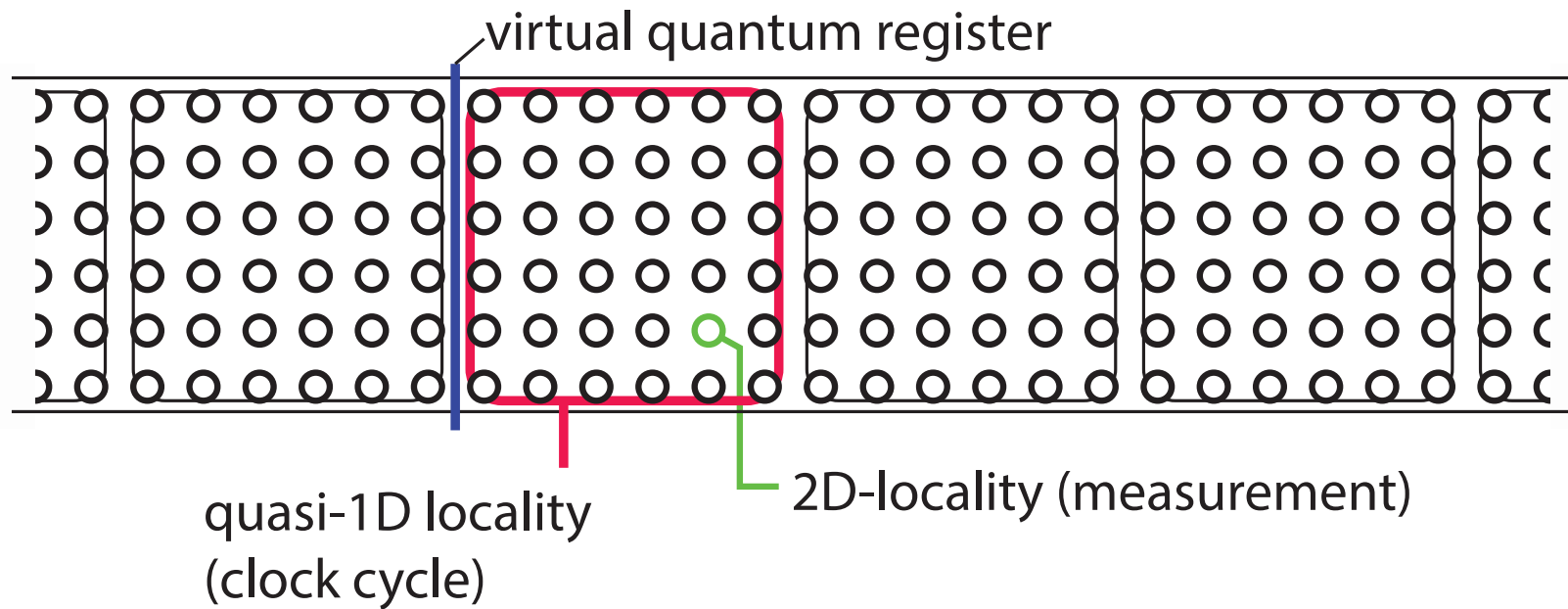
The clock cycle:



- Every byproduct operator is mapped back to itself after n columns ($n = \text{circumference}$).

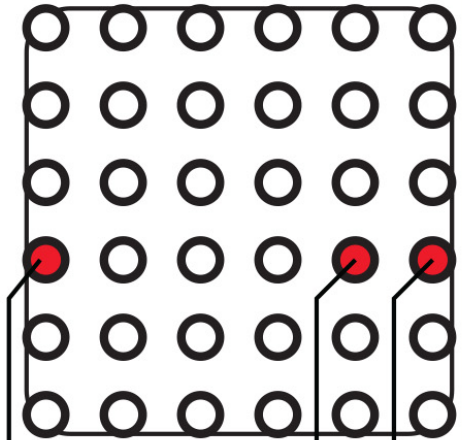
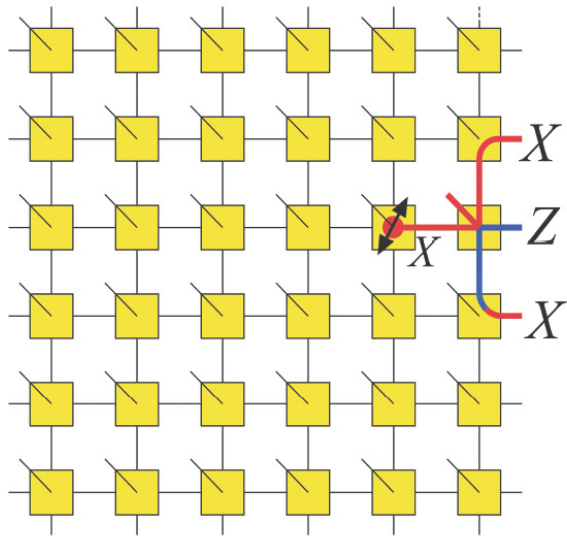
\Rightarrow If a gate can be done once, it can be done many times.

B: Cluster-like \Rightarrow universal



- Map 2D system to effective 1D system

B: Cluster-like \Rightarrow universal



$$e^{i\alpha Z_k}$$

$$e^{i\alpha X_{k-1}Z_kX_{k+1}}$$

$$e^{i\alpha X_k}$$

Universal gate set on $n/2$ qubits

B: Cluster-like \Rightarrow universal

2D cluster state:



$$e^{i d \alpha Z_k}$$

$$e^{i d \alpha X_{k-1} Z_k X_{k+1}}$$

$$e^{i d \alpha X_k}$$

Throughout the phase:

$$e^{i |\nu| d \alpha Z_k}$$

$$e^{i |\nu| d \alpha X_{k-1} Z_k X_{k+1}}$$

$$e^{i |\nu| d \alpha X_k}$$

$$|\nu| \leq 1$$

(ν depends on the location in the phase)

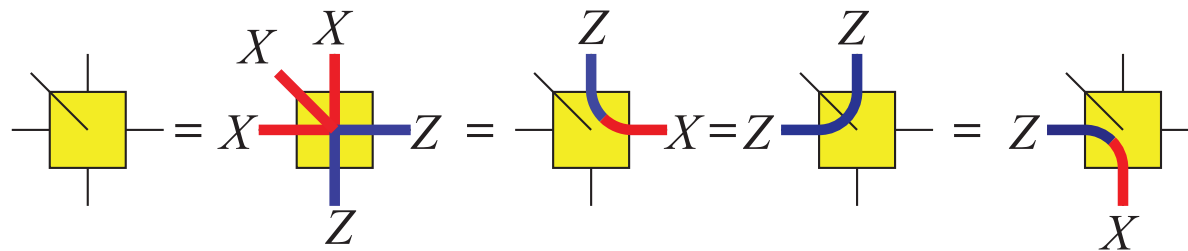
About ν : RR, A.Prakash, D.-S. Wang, T.-C.Wei, D.T. Stephen, Phys. Rev. A (2017).

Summary and outlook

- There exists a symmetry-protected phase in 2D with uniform universal computational power for MBQC.
- *Can we have a classification of MBQC schemes in 2D, based on symmetry?*
- Symmetry Lego is fun—Try it!

PRL 122, 090501 (2019)

Related: **Quantum 3, 162 (2019)**



A: In cluster phase \Rightarrow cluster-like

Lemma 3. [*] Symmetric gapped ground states in the same SPT phase are connected by symmetric local quantum circuits of constant depth.

For any state $|\Phi\rangle$ in the phase,

$$|\Phi\rangle = U_k U_{k-1} \dots U_1 |2D \text{ cluster}\rangle.$$

Look at an individual symmetry-respecting gate in the circuit,

$$U = \sum_j c_j T_j, \text{ with } T_j \in \mathcal{P}.$$

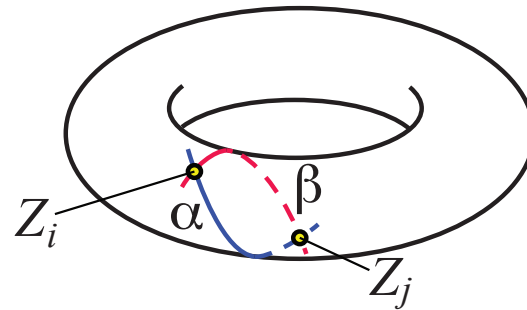
Which Pauli observables T_j can be admitted in the expansion?

[*] X. Chen, Z.C. Gu, and X.G. Wen, Phys. Rev. B **82**, 155138 (2010).

A: In cluster phase \Rightarrow cluster-like

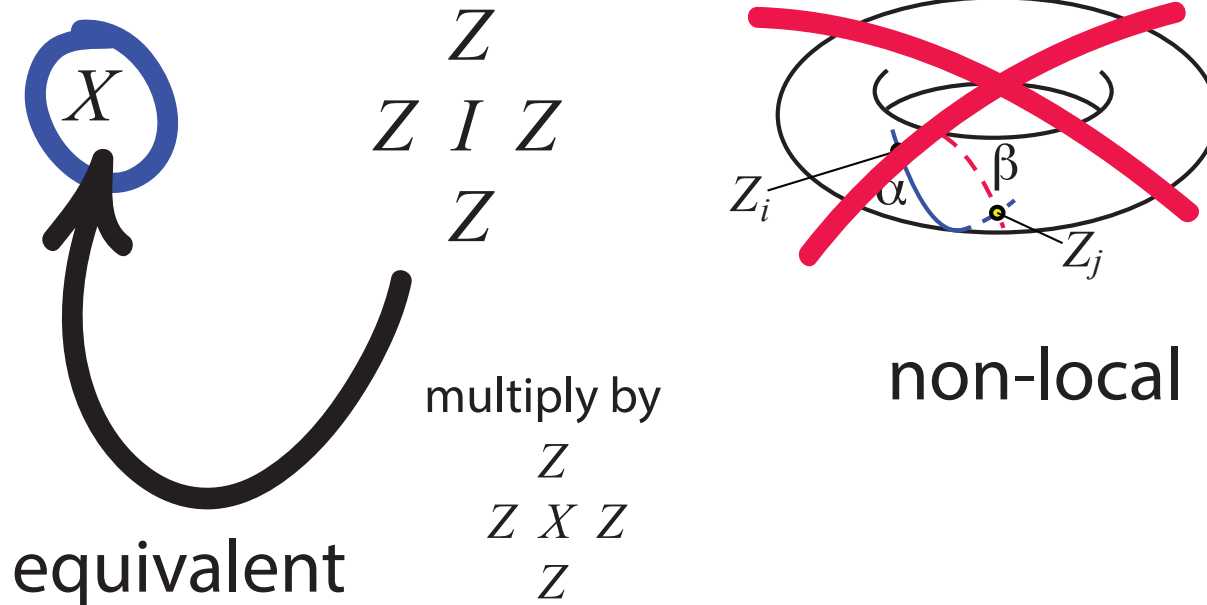
Which Paulis T_j can be admitted in the expansion $U = \sum_j c_j T_j$?

$$X \quad \begin{matrix} Z \\ Z I Z \\ Z \end{matrix}$$



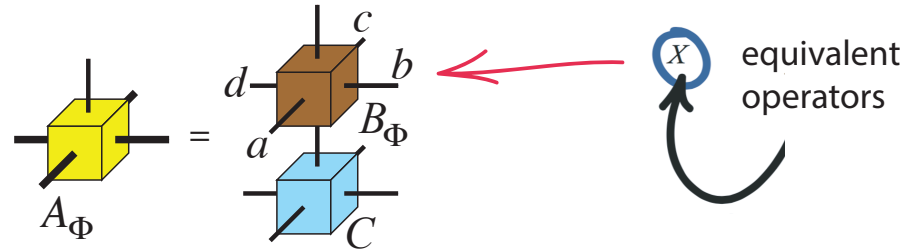
A: In cluster phase \Rightarrow cluster-like

Which Paulis T_j can be admitted in the expansion $U = \sum_j c_j T_j$?

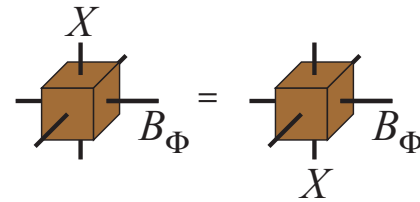


Only X -type Pauli operators survive in the expansion.

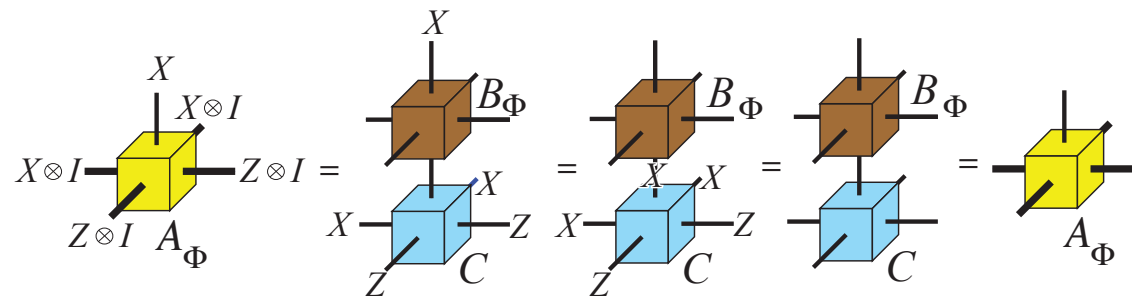
Description of the local tensors:



With



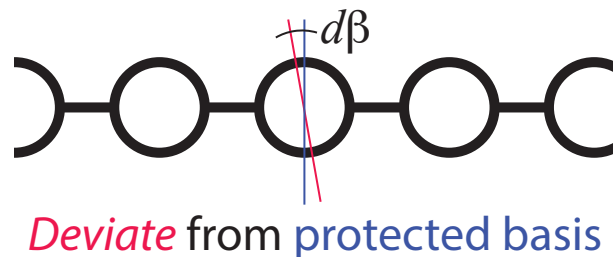
Hence



- Local tensors A_Φ describing $|\Phi\rangle$ are invariant under the cluster-like symmetries.

The parameter ν

There is a complex-valued parameter ν , $|\nu| \leq 1$, that needs to be known about the location of the resource state within the phase.



For infinitesimal angles $d\beta$, this results in a logical rotation [*]

$$e^{id\beta|\nu|T},$$

for some Pauli operator T . (E.g., $T = Z_k, X_k, X_{k-1}Z_kX_{k+1}$).

We require that $\nu \neq 0$.

[*] RR, D.-S. Wang, A. Prakash, T.-C. Wei, D.T. Stephen, PRA 96 (2017).